These are the examples for Tutorial 2 with solutions. The alphabet is $\Sigma=\{0,1\}$ in all the examples.

## Problem 1

$L_{1}=\left\{0^{n^{2}} \mid n\right.$ is an integer and $\left.n \geq 0\right\}$

## Solution

Suppose it is regular, then there exist a DFA that accept $L_{1}$ and there are $N$ states in the minimal DFA, say $S$. Let us choose the string $z=0^{N^{2}}$. By pumping lemma, $z=u v w$, where $|u v| \leq N$ and $|v| \geq 1, u v^{i} w$ can be accepted by $S$ for every integer $i \geq 0$. Especially, when $i=2, u v^{2} w$ can be accepted by $S$. Let us check whether $u v^{2} w$ is in $L_{1}$.

$$
N^{2}=|u v w|<\left|u v^{2} w\right|=|u v w|+|v| \leq|u v w|+|u v| \leq N^{2}+N<N^{2}+2 N+1=(N+1)^{2}
$$

The length of $u v^{2} w$ is not square of any integers, then $u v^{2} w$ is not in $L_{1}$, contradiction.

## Problem 2

$L_{2}=\left\{0^{m} 1^{n} \mid m>n \geq 0\right\}$

## Solution

Suppose $L_{2}$ is regular, then there exist a DFA, say $S$, with $N$ states that can accept $L_{2}$. Choose $z=0^{N} 1^{N-1}$ and $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. With only pumping deleted, $u w$ also can be accepted by $S$. Notice that all symbols in $v$ are 0 s, then $u w \notin L_{2}$, contradiction.

## Problem 3

$L_{3}=\left\{0^{2 n} \mid n \geq 1\right\}$

## Solution

It is easy to construct a DFA for $L_{3}$, so it is regular.

## Problem 4

$L_{4}=\left\{0^{m} 1^{n} 0^{m+n} \mid m \geq 1\right.$ and $\left.n \geq 1\right\}$

## Solution

Suppose $L_{4}$ is regular, then there is a DFA, say $S$, with $N$ states that can accept $L_{4}$. Choose $z=0^{N} 1^{N} 0^{2 N}$ and $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. With only pumping deleted, $u w$ also can be accepted by $S$. Notice that all symbols in $v$ are 0 s , then $u w \notin L_{4}$, contradiction.

## Problem 5

$L_{5}=\left\{0^{n} \mid n\right.$ is a prime $\}$

## Solution

Suppose it is regular, then there is a DFA, say $S$, with $N$ states that can accept $L_{5}$. Choose $z=0^{p}$, where $p$ is a prime and $p \geq N$. We have the patition $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. Consider the string $z^{\prime}=u v^{|z|+1} w$, it can be accepted by $L_{5}$, but $\left|z^{\prime}\right|=|z|+|z||v|=|z|(1+|v|)$ is not a prime, contradiction.

## Problem 6

$L_{6}=\{x \mid x$ does not have three consecutive 0s $\}$.

## Solution

$L_{6}$ is regular. You can construct a DFA for it.

## Problem 7

$L_{7}=\{x \mid x$ has an equal number of 0 s and 1 s$\}$.

## Solution

Suppose $L_{7}$ is regular, then there is a DFA, say $S$, with $N$ states that can accept $L_{7}$. Choose $z=0^{N} 1^{N}$ and $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. With only pumping deleted, $u w$ also can be accepted by $S$. Notice that all symbols in $v$ are 0 s, then there are less 0's than 1's in $u w, u w \notin L_{7}$, contradiction.

## Problem 8

$L_{8}=\left\{x \mid x=x^{R}\right\}$. Recall that $x^{R}$ is $x$ written backwards; for example, $(011)^{R}=110$.

## Solution

Suppose $L_{8}$ is regular, then there is a DFA, say $S$, with $N$ states that can accept $L_{8}$. Choose $z=0^{N} 110^{N}$ and $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. With only pumping deleted, $u w$ also can be accepted by $S$. Notice that all symbols in $v$ are 0 s, then $u w \notin L_{8}$, contradiction.

## Problem 9

$L_{9}=\{x \mid x$ has a different number of 0s and 1 s$\}$.

## Solution

The easier way to prove $L_{9}$ is not regular goes like this. Suppose it is regular, then $L_{7}$ is $L_{9}$ 's complement, hence $L_{7}$ is regular, contradiction.

If you want to prove $L_{9}$ is not regular using the pumping lemma, it is also possible, but a bit more difficult. Suppose it is regular, then there is a DFA, say $S$, with $N$ states that can accept $L_{9}$. Choose $z=0^{N} 1^{N+N!}$ and $z=u v w$, where $|u v| \leq N$ and $v$ is the pumping, thus $|v| \geq 1$. (Here $N!=1 \cdot 2 \cdots N$.) Then $z^{\prime}=u v^{i} w$ can be accepted by $S$ for every nonnegtive integer $i$. Set $i=N!/|v|+1$, then $z^{\prime}=u v^{N!/|v|+1} w$. Notice that all symbols in $v$ are 0 s, it is easy to check $z^{\prime} \notin L_{9}$, contradiction.

